

# Midterm 2 Solutions

-1-

FCE 564/645

Spring, 2013

1) (a) This is a  $(6, 2)$  code:

- $r = k/n = 2/6 = 1/3$

- $\mathcal{C} = \{000000, 101010, 010101, 111111\}$

- $A(x) = 1 + 2x^3 + x^6, d_{\min} = 3$

- $G = [I_2 | A] \Rightarrow H = [A^T | I_4]$

$$= \begin{bmatrix} 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 \end{bmatrix}$$

- $2^{n-k} = 16$

(b)

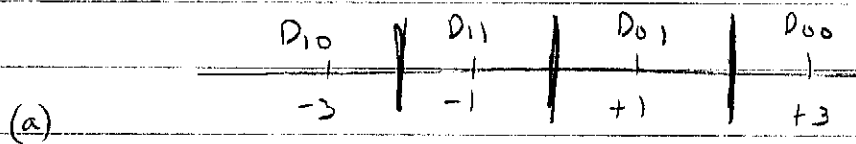
No

$$\mathcal{C} = \{000000, 111100, 001111, 111111\}$$

$$G = \begin{bmatrix} 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 \end{bmatrix}$$

has  $d_{\min} = 4$ .

2)



(a)

$-1.0, -2.1, -2.1, 2.1, 1.2, 0.3$   
 $\downarrow \quad \downarrow \quad \downarrow \quad \downarrow \quad \downarrow \quad \downarrow$   
 $11 \quad 10 \quad 10 \quad 00 \quad 01 \quad 01$

Hamming distance

→ closest to  
101010

closest to  
010101

↓  
 $\hat{a} = 10$

↓  
 $\hat{a} = 01$

output 1001

(b) think

$$00 \Rightarrow 000000 \Rightarrow (+3, +3, +3)$$

$$10 \Rightarrow 101010 \Rightarrow (-3, -3, -3)$$

$$01 \Rightarrow 010101 \Rightarrow (+1, +1, +1)$$

$$11 \Rightarrow 111111 \Rightarrow (-1, -1, -1)$$

Euclidean distance

$-1.0, -2.1, -2.1, 1.5, 1.9, 3.5$

closest to  
 $(-1, -1, -1)$

closest to  
 $(+3, +3, +3)$

11

00

for example:

$$d_E^2(y, s_{00}) = (-1-3)^2 + (-2.1-3)^2 + (-2.1-3)^2$$

c)

$$P(E) \leq \sum_{i=1}^3 P(\text{choose } c_i | c_0, s_{k+1})$$

$$= \sum_{i=1}^3 Q\left(\frac{d_E(c_i, c_0)}{\sqrt{2N_0}}\right)$$

$$= Q\left(\sqrt{\frac{108}{2N_0}}\right) + Q\left(\sqrt{\frac{12}{2N_0}}\right) + Q\left(\sqrt{\frac{48}{2N_0}}\right)$$

$$= Q\left(\sqrt{\frac{6}{N_0}}\right) + Q\left(\sqrt{\frac{24}{N_0}}\right) + Q\left(\sqrt{\frac{54}{N_0}}\right)$$

3)

(a)

Need to find three codewords of wt.  $\geq 5$   
that are still a distance 5 apart.  
But any two wt. 5 codewords overlap  
in at least 3 places:

1111100  
0011111

and thus  $d_{\min} = 4$  for a  $(7,2)$  code.

(b)

• Start with an  $(n, k)$  code:  $r = n - k$

Shortened, it is  $(n-1, k-1)$   $r = \frac{k-1}{n-1} = \frac{k}{n} - \frac{1}{n-1}$

Smaller

• Recall

$d_{\min} = \min \#$  of columns that sum to  $Q$

removing columns of  $H$  can leave this  
the same or make it larger.

(c)

• All columns are distinct  $\Rightarrow d_{\min} \geq 2$

$$\text{but } \begin{pmatrix} 0 \\ 0 \\ 0 \\ 0 \\ 1 \end{pmatrix} + \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \\ 0 \end{pmatrix}$$

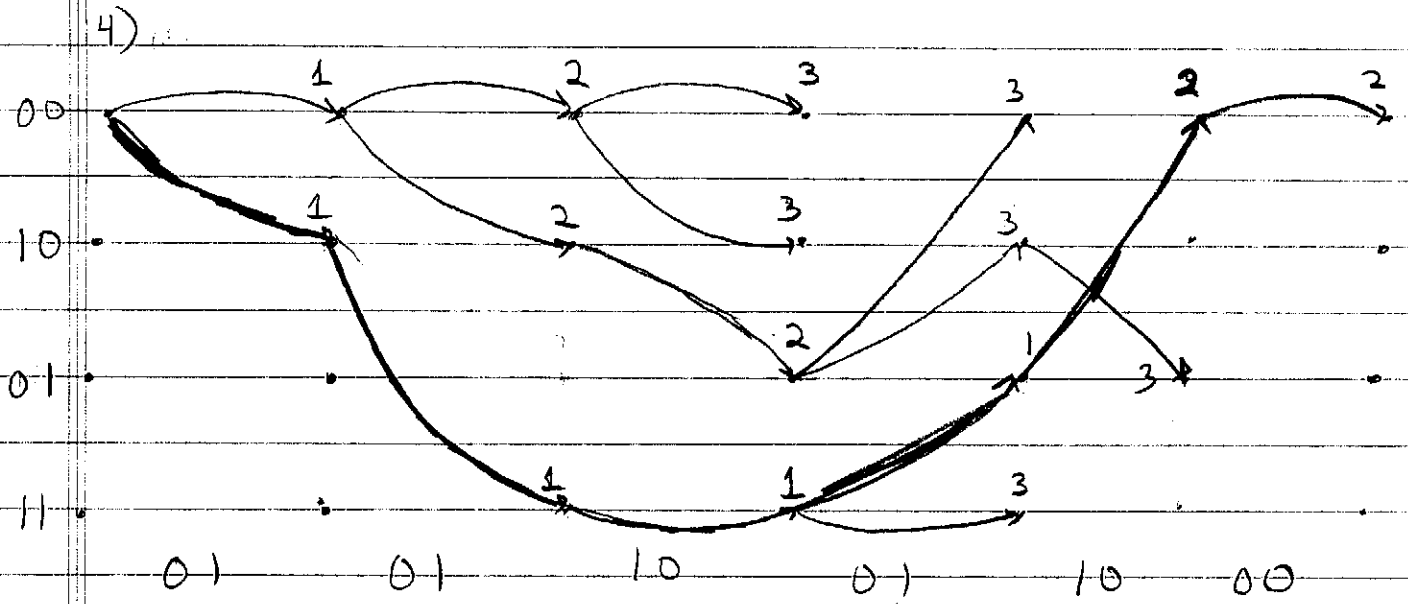
$$\Rightarrow d_{\min} \leq 3; \text{ thus, } d_{\min} = 3$$

•  $2^{n-k} = 2^5 = 32$ , which is the right number for the all 0's error pattern and every wt 1 pattern (of which there are 31).

• There is no (7,2),  $d_{\min} = 5$  code; thus, by part (b), the best I can get is the (6,1) repetition code,

$$C = \{000000, 111111\}$$

$$H = \begin{bmatrix} 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$



Output

1      1      1      0

4

7  
5

5) (a)

$$\tilde{G} = H$$

$$\Rightarrow \tilde{G} \rightarrow (n-k) \times n \Rightarrow \text{length} = n$$

$$\text{bits codeword} = n-k$$

$$\text{rate} = \frac{n-k}{n} = 1 - k/n = 1 - r$$

(b)

$$x \oplus c_i = u^T \tilde{G} c_i^T = u^T H c_i^T = 0$$

$$\text{since } H c_i^T = 0$$

$$(c) \text{ No. } 2^k + 2^{n-k} < 2^n$$

$$(d) \tilde{G} = \begin{pmatrix} 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 & 0 & 1 \end{pmatrix}$$

$$c = \{0000000, 0001111, 0110011, 1010101,$$

$$0111100, 1011010, 1100110, 1101001\}$$

$$A(x) = 1 + 7x^4$$