

Final Exam Solutions

- 1 -

FCE 564/645

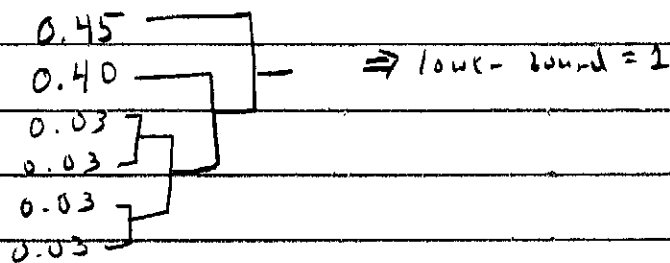
Spring 2019

1)(a)

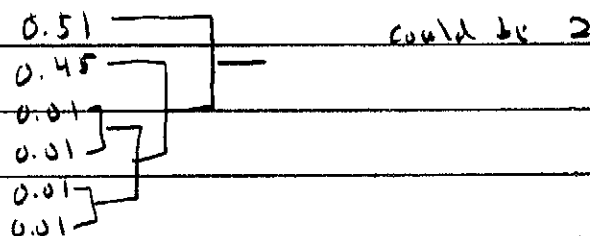
| | | | <u>code</u> | <u>p(x) l(x)</u> |
|----|------|--|-------------|------------------|
| AA | 0.16 | | 000 | 0.48 |
| AB | 0.16 | | 001 | 0.48 |
| AC | 0.08 | | 110 | 0.24 |
| BA | 0.16 | | 10 | 0.32 |
| BB | 0.16 | | 010 | 0.48 |
| BC | 0.08 | | 0110 | 0.32 |
| CA | 0.08 | | 0111 | 0.32 |
| CB | 0.08 | | 1110 | 0.32 |
| CC | 0.04 | | 1111 | 0.16 |
| | | | | 3.12 |

$$R = 3.12 / 2 = 1.56 \text{ bits/symbol}$$

Certainly, you could have a 1-bit sequence



Could it be larger?



But, being the second largest - and - all others (except largest) sum to $\leq 0.45 \Rightarrow l=2$ so $\max 1 \leq l(x) \leq 2$

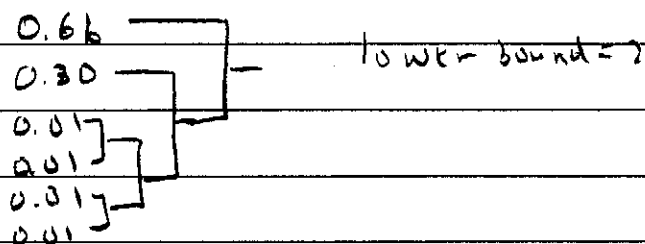
Not on exam (checked)

(b) Can it be $l=1$?

No.

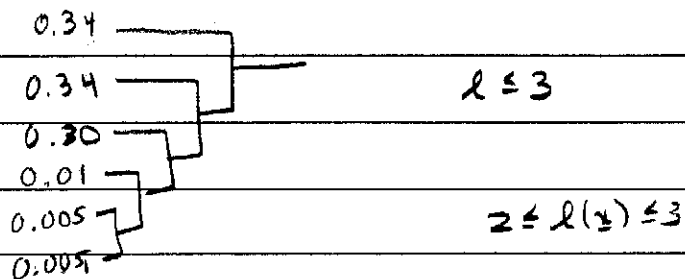
When there are three blocks left, 0.30 cannot be the largest.

Can I have $l=2$? yes



Upper bound?

At most the third largest block



$$2) (a) \int_0^1 \cos(2\pi t) dt = \frac{1}{2\pi} \sin(2\pi t) \Big|_0^1 = 0 - 0 = 0$$

$$\int_0^1 \sin(2\pi t) dt = \frac{-1}{2\pi} \cos(2\pi t) \Big|_0^1 = -1 + 1 = 0$$

$$\int_0^1 \sin(2\pi t) \cos(2\pi t) dt = \int_0^1 \frac{1}{2} \sin(4\pi t) dt = \frac{-1}{8\pi} \cos(4\pi t) \Big|_0^1 = -1 + 1 = 0$$

$$\int_0^1 \cos^2(2\pi t) dt = \int_0^1 \left(\frac{1}{2} + \frac{1}{2} \cos(4\pi t) \right) dt$$

$$= \frac{1}{2} + \frac{1}{4\pi} \sin(4\pi t) \Big|_0^1 = \frac{1}{2}$$

(b) From part (a), 1, $\cos(2\pi t)$, and $\sin(2\pi t)$ form an orthogonal set. Thus, $N \geq 3$.

Try:

$$\phi_1(t) = 1, \quad 0 \leq t \leq 1$$

$$\phi_2(t) = \sqrt{2} \cos(2\pi t), \quad 0 \leq t \leq 1$$

$$\phi_3(t) = \sqrt{2} \sin(2\pi t), \quad 0 \leq t \leq 1$$

All four signals can be written in terms of $\phi_1(t), \phi_2(t), \phi_3(t)$.

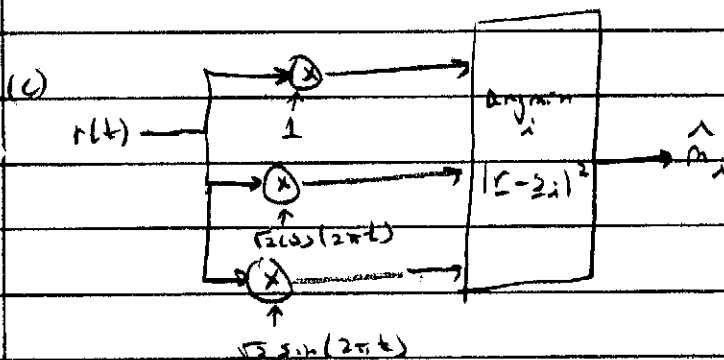
$$s_1 = (2, 0, 0)^T$$

$$s_2 = (1, 2, 0)^T$$

$$s_3 = (1, 0, 2)^T$$

$$s_4 = (-1, 0, 2)^T$$

Also, could do G-S easily, using (a)



(d)

$$P(E) \leq \frac{1}{m} \sum_{i=1}^m \sum_{j=1}^m Q\left(\frac{d_{ij}}{\sqrt{2N_0}}\right)$$

| | 1 | 2 | 3 | 4 |
|---|-------------|-------------|------------|-------------|
| 1 | - | $\sqrt{5}$ | $\sqrt{5}$ | $\sqrt{13}$ |
| 2 | $\sqrt{5}$ | - | $\sqrt{8}$ | $\sqrt{12}$ |
| 3 | $\sqrt{5}$ | $\sqrt{8}$ | - | 2 |
| 4 | $\sqrt{13}$ | $\sqrt{12}$ | 2 | - |

$$\leq \frac{1}{4} \left(4Q\left(\sqrt{\frac{5}{2N_0}}\right) + 2Q\left(\sqrt{\frac{4}{N_0}}\right) + 2Q\left(\sqrt{\frac{13}{2N_0}}\right) + 2Q\left(\sqrt{\frac{6}{N_0}}\right) + 2Q\left(\sqrt{\frac{2}{N_0}}\right) \right)$$

$$E_s = \frac{1}{4} (4 + 5 + 5 + 5)$$

$$= 19/4$$

dominant term

$$\frac{1}{2} Q\left(\sqrt{\frac{2}{N_0}}\right) = \frac{1}{2} Q\left(\sqrt{\frac{2 \cdot 4/19 \cdot 19/4}{N_0}}\right) = \frac{1}{2} Q\left(\sqrt{\frac{8E_s}{19N_0}}\right)$$

[and

$$E_s = E_b \log_2 m = 2E_b]$$

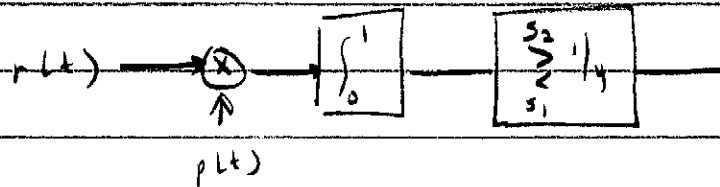
$$= \frac{1}{2} Q\left(\sqrt{\frac{16E_b}{19N_0}}\right)$$

3) (a)

$$\phi(t) = p(t)$$

$$s_1 = -1/4$$

$$s_2 = 3/4$$



$$(b) P_b = Q\left(\frac{0}{\sqrt{2N_0}}\right) = Q\left(\frac{1}{\sqrt{2N_0}}\right) = Q\left(\frac{\sqrt{10/5} \cdot \sqrt{5/16}}{2N_0}\right) = Q\left(\frac{8}{\sqrt{5}} \cdot \frac{E_b}{2N_0}\right)$$

$$E_b = \frac{1}{2} \left(\left(\frac{1}{4}\right)^2 + \left(\frac{3}{4}\right)^2 \right) = \frac{5}{16}$$

$$(c) 10 \log_{10} \left(\frac{2}{3/5} \right) = 10 \log_{10} 5/3 \text{ dB worse!}$$

(d) Need three linearly independent codewords:

$$G = \begin{bmatrix} 1 & 0 & 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 \end{bmatrix} \quad \left. \vphantom{\begin{bmatrix} 1 & 0 & 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 \end{bmatrix}} \right\} k=3 \quad r = \frac{k}{n} = \frac{3}{6} = \frac{1}{2}$$

$$C = \{000000, 001110, 110010, 111100, 101001, 100111, 011011, 010101\}$$

$$A(x) = 1 + 4x^3 + 3x^4$$

| \underline{s} | leader |
|-----------------|--------|
| 000 | 000000 |
| 001 | 000001 |
| 010 | 000010 |
| 011 | 000011 |
| 100 | 001000 |
| 101 | 100000 |
| 110 | 000100 |
| 111 | 010000 |

$$(e) \quad p = Q\left(\sqrt{\frac{100}{5}} \frac{E_s}{N_0}\right) = Q\left(\sqrt{\frac{4}{5}} \frac{E_b}{N_0}\right)$$

$$r \frac{E_b}{N_0} = \frac{E_s}{N_0}$$

$$E_b = 2E_s$$

$$\text{and } P(E) = 1 - P(L) = 1 - (1-p)^6 - 6p(1-p)^5 - p^2(1-p)^4$$

$$(f) \quad P(E) \leq \sum_{i=0}^3 Q\left(\frac{A_{0,i}}{\sqrt{2}N_0}\right)$$

$$= 4Q\left(\sqrt{\frac{3 \cdot 1^2}{2N_0}}\right) + 3Q\left(\sqrt{\frac{4 \cdot 1^2}{2N_0}}\right)$$

$$= 4Q\left(\sqrt{\frac{3 \cdot 16 \cdot 5 \cdot 5}{16 \cdot 2N_0}}\right) + 3Q\left(\sqrt{\frac{4 \cdot 16 \cdot 5 \cdot 5}{16 \cdot 2N_0}}\right)$$

$$= 4Q\left(\sqrt{\frac{24}{5}} \frac{E_s}{N_0}\right) + 3Q\left(\sqrt{\frac{32}{5}} \frac{E_s}{N_0}\right)$$

$$[E_b = 2E_s]$$

$$= 4Q\left(\sqrt{\frac{12}{5}} \frac{E_b}{N_0}\right) + 3Q\left(\sqrt{\frac{16}{5}} \frac{E_b}{N_0}\right)$$

4) (a)

For a t -error correcting code, there are

$$1 + \binom{n}{1} + \binom{n}{2} + \dots + \binom{n}{t}$$

patterns to be corrected. Each must have its own syndrome; thus,

codewords

$$2^{n-k} \geq 1 + \binom{n}{1} + \binom{n}{2} + \dots + \binom{n}{t}$$
$$\Rightarrow 2^k \leq 2^n / (1 + \binom{n}{1} + \binom{n}{2} + \dots + \binom{n}{t})$$

(b)

$$H = \begin{bmatrix} 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 & 0 & 1 \end{bmatrix}$$

$$k=4, n=7, t=1$$

$$3 = \log_2(1+7) = 3 \quad \checkmark$$